

Computer Vision II - Lecture 13

Multi-Object Tracking III

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Course Outline

- Single-Object Tracking
- Bayesian Filtering
 - Kalman filters
 - Particle filters
 - Case studies
- Multi-Object Tracking
 - Introduction
 - MHT, JPDAF
 - Network Flow Optimization
- Articulated Tracking



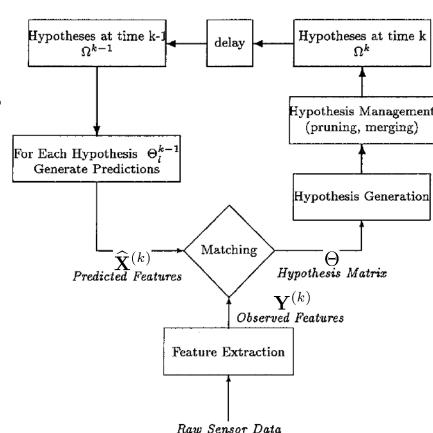
Topics of This Lecture

- Recap: MHT
- Data Association as Linear Assignment Problem
 - LAP formulation
 - Greedy algorithm
 - > Hungarian algorithm
- Tracking as Network Flow Optimization
 - Min-cost network flow
 - Generalizing to multiple frames
 - Complications
 - > Formulation

Recap: Multi-Hypothesis Tracking (MHT)

Ideas

- Instead of forming a track tree, keep a set of hypotheses that generate child hypotheses based on the associations.
- Enforce exclusion constraints between tracks and measurements in the assignment.
- Integrate track generation into the assignment process.
- After hypothesis generation, merge and prune the current hypothesis set.



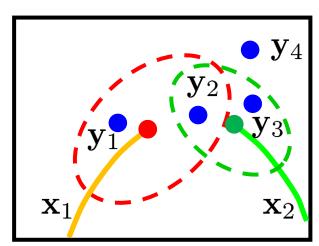
D. Reid, <u>An Algorithm for Tracking Multiple Targets</u>, IEEE Trans. Automatic Control, Vol. 24(6), pp. 843-854, 1979.



Recap: Hypothesis Generation

Create hypothesis matrix of the feasible associations

$$egin{aligned} \mathbf{x}_1 & \mathbf{x}_2 \mathbf{x}_{fa} \mathbf{x}_{nt} \\ \Theta &= egin{bmatrix} 1 & 0 & 1 & 1 & \mathbf{y}_1 \\ 1 & 1 & 1 & 1 & \mathbf{y}_2 \\ 0 & 1 & 1 & 1 & \mathbf{y}_3 \\ 0 & 0 & 1 & 1 & \mathbf{y}_4 \end{bmatrix} \end{aligned}$$



- Interpretation
 - > Columns represent tracked objects, rows encode measurements
 - A non-zero element at matrix position (i,j) denotes that measurement \mathbf{y}_i is contained in the validation region of track \mathbf{x}_j .
 - ightharpoonup Extra column \mathbf{x}_{fa} for association as false alarm.
 - ightharpoonup Extra column \mathbf{x}_{nt} for association as *new track*.
 - Turn this hypothesis matrix



Recap: Assignments

| $_ Z_j$ | \mathbf{x}_1 | \mathbf{x}_2 | \mathbf{x}_{fa} | \mathbf{x}_{nt} |
|----------------|----------------|----------------|-------------------|-------------------|
| \mathbf{y}_1 | 0 | 0 | 1 | 0 |
| \mathbf{y}_2 | 1 | 0 | 0 | 0 |
| \mathbf{y}_3 | 0 | 1 | 0 | 0 |
| \mathbf{y}_4 | 0 | 0 | 0 | 1 |

Impose constraints

- A measurement can originate from only one object.
- ⇒ Any row has only a single non-zero value.
- An object can have at most one associated measurement per time step.
- \Rightarrow Any column has only a single non-zero value, except for \mathbf{x}_{fa} , \mathbf{x}_{nt}

Recap: Calculating Hypothesis Probabilities

Probabilistic formulation

- It is straightforward to enumerate all possible assignments.
- However, we also need to calculate the probability of each child hypothesis.
- This is done recursively:

$$\begin{split} p(\Omega_j^{(k)}|\mathbf{Y}^{(k)}) &= p(Z_j^{(k)}, \Omega_{p(j)}^{(k-1)}|\mathbf{Y}^{(k)}) \\ &\stackrel{Bayes}{=} \eta p(\mathbf{Y}^{(k)}|Z_j^{(k)}, \Omega_{p(j)}^{(k-1)}) p(Z_j^{(k)}, \Omega_{p(j)}^{(k-1)}) \\ &= \eta p(\mathbf{Y}^{(k)}|Z_j^{(k)}, \Omega_{p(j)}^{(k-1)}) p(Z_j^{(k)}|\Omega_{p(j)}^{(k-1)}) p(\Omega_{p(j)}^{(k-1)}) \\ &\stackrel{\text{Normalization}}{\text{factor}} \quad \text{Measurement} \quad \underset{\text{likelihood}}{\text{Prob. of}} \quad \underset{\text{assignment set}}{\text{Prob. of}} \quad \underset{\text{parent}}{\text{Prob. of}} \end{split}$$



Recap: Measurement Likelihood

Use KF prediction

- Assume that a measurement $\mathbf{y}_i^{(k)}$ associated to a track \mathbf{x}_j has a Gaussian pdf centered around the measurement prediction $\hat{\mathbf{x}}_j^{(k)}$ with innovation covariance $\widehat{\boldsymbol{\Sigma}}_j^{(k)}$.
- Further assume that the pdf of a measurement belonging to a new track or false alarm is uniform in the observation volume W (the sensor's field-of-view) with probability $W^{\text{-}1}$.
- Thus, the measurement likelihood can be expressed as

$$p\left(\mathbf{Y}^{(k)}|Z_{j}^{(k)},\Omega_{p(j)}^{(k-1)}\right) = \prod_{i=1}^{M_{k}} \mathcal{N}\left(\mathbf{y}_{i}^{(k)};\hat{\mathbf{x}}_{j},\widehat{\boldsymbol{\Sigma}}_{j}^{(k)}\right)^{\delta_{i}} W^{-(1-\delta_{i})}$$

$$= W^{-(N_{fal}+N_{new})} \prod_{i=1}^{M_{k}} \mathcal{N}\left(\mathbf{y}_{i}^{(k)};\hat{\mathbf{x}}_{j},\widehat{\boldsymbol{\Sigma}}_{j}^{(k)}\right)^{\delta_{i}}$$

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Recap: Probability of an Assignment Set

$$p(Z_j^{(k)}|\Omega_{p(j)}^{(k-1)})$$

- Composed of three terms
 - 1. Probability of the number of tracks N_{det} , N_{fal} , N_{new}
 - Assumption 1: N_{det} follows a binomial distribution

$$p(N_{det}|\Omega_{p(j)}^{(k-1)}) = {N \choose N_{det}} p_{det}^{N_{det}} (1 - p_{det})^{(N-N_{det})}$$

where N is the number of tracks in the parent hypothesis

- Assumption 2: N_{fal} and N_{new} both follow a Poisson distribution with expected number of events $\lambda_{fal}W$ and $\lambda_{new}W$

$$p(N_{det}, N_{fal}, N_{new} | \Omega_{p(j)}^{(k-1)}) = {N \choose N_{det}} p_{det}^{N_{det}} (1 - p_{det})^{(N-N_{det})} \cdot \mu(N_{fal}; \lambda_{fal}W) \cdot \mu(N_{new}; \lambda_{new}W)$$

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Recap: Probability of an Assignment Set

2. Probability of a specific assignment of measurements

- Such that $M_k = N_{det} + N_{fal} + N_{new}$ holds.
- This is determined as 1 over the number of combinations

$$\begin{pmatrix} M_k \\ N_{det} \end{pmatrix} \begin{pmatrix} M_k - N_{det} \\ N_{fal} \end{pmatrix} \begin{pmatrix} M_k - N_{det} - N_{fal} \\ N_{new} \end{pmatrix}$$

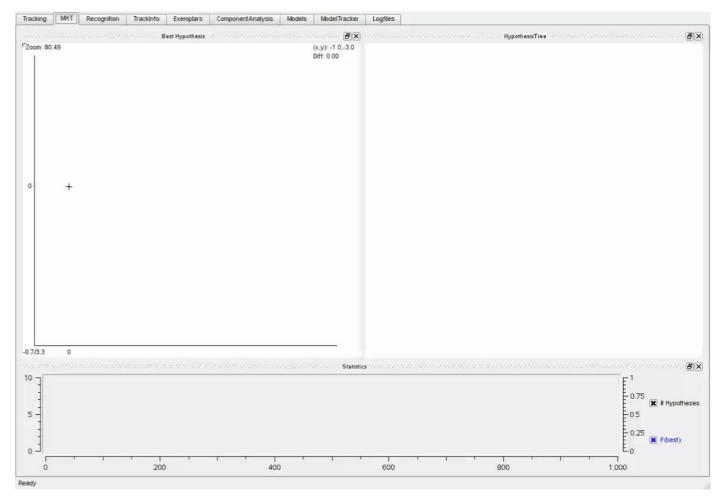
3. Probability of a specific assignment of tracks

- Given that a track can be either detected or not detected.
- This is determined as 1 over the number of assignments

$$\frac{N!}{(N-N_{det})!} \left(\begin{array}{c} N-N_{det} \\ N_{det} \end{array} \right)$$

⇒ When combining the different parts, many terms cancel out!

Laser-based Leg Tracking using MHT



K. Arras, S. Grzonka, M. Luber, W. Burgard, Efficient People Tracking in Laser Range Data using a Multi-Hypothesis Leg-Tracker with Adaptive Occlusion Probabilities, ICRA'08.

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Laser-based People Tracking using MHT

Multi Hypothesis Tracking of People

Matthias Luber, Gian Diego Tipaldi and Kai O. Arras

Laser-baser People Tracking using MHT (Inner city of Freiburg, Germany) Results projected onto video data.







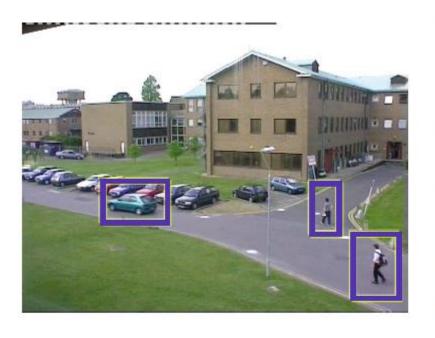
Topics of This Lecture

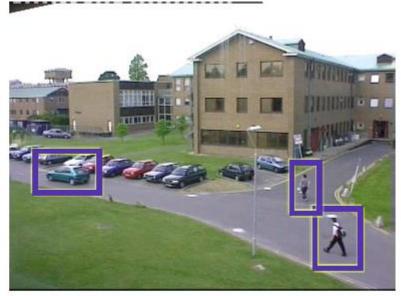
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Back to Data Association...

Goal: Match detections across frames











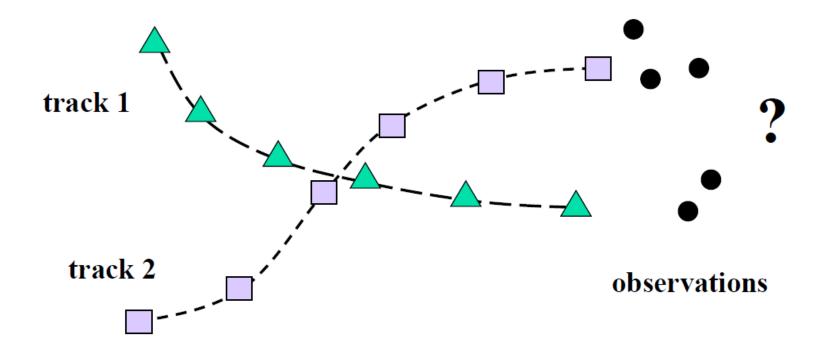








Data Association



- Main question here
 - How to determine which measurements to add to which track?
 - Today: consider this as a matching problem



Linear Assignment Formulation

Form a matrix of pairwise similarity scores

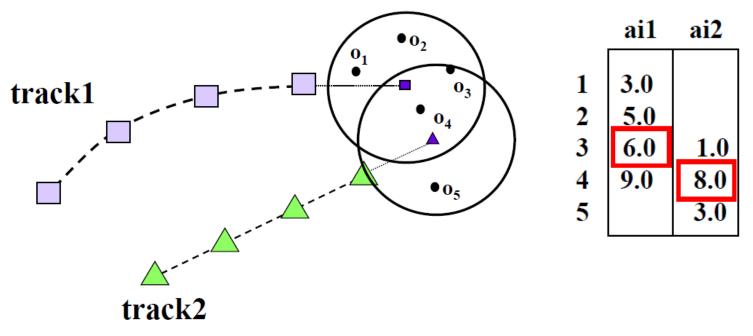
| | | • | Frame $t+1$ | I |
|--|-----|------|-------------|------|
| Similarity could bebased on motion prediction | | X | | 4 |
| based on appearancebased on both | | 0.11 | 0.95 | 0.23 |
| Frame t | A | 0.85 | 0.25 | 0.89 |
| • Goal | THE | 0.90 | 0.12 | 0.81 |

- - Choose one match from each row and column to maximize the sum of scores



Linear Assignment Formulation

- Example: Similarity based on motion prediction
 - Predict motion for each trajectory and assign scores for each measurement based on inverse (Mahalanobis) distance, such that closer measurements get higher scores.



Choose at most one match in each row and column to maximize sum of scores



Linear Assignment Problem

- Formal definition
 - > Maximize $\sum \sum w_{ij}z_{ij}$

subject to
$$\sum_{j=1} z_{ij}=1;\;i=1,2,\ldots,N$$
 $\sum_{i=1} z_{ij}=1;\;j=1,2,\ldots,M$ Those constraints ensure that Z is a permutation matrix

- The permutation matrix constraint ensures that we can only match up one object from each row and column.
- Note: Alternatively, we can minimize cost rather than maximizing weights.

$$\arg\min_{z_{ij}} \sum_{i=1}^{N} \sum_{j=1}^{M} c_{ij} z_{ij}$$



Greedy Solution to LAP

| | 1 | 2 | 3 | 4 | 5 |
|---|------|------|------|------|------|
| 1 | 0.95 | 0.76 | 0.62 | 0.41 | 0.06 |
| 2 | 0.23 | 0.46 | 0.79 | 0.94 | 0.35 |
| 3 | 0.61 | 0.02 | 0.92 | 0.92 | 0.81 |
| 4 | 0.49 | 0.82 | 0.74 | 0.41 | 0.01 |
| 5 | 0.89 | 0.44 | 0.18 | 0.89 | 0.14 |

Greedy algorithm

- > Find the largest score
- Remove scores in same row and column from consideration
- Repeat
- Result: score =



Greedy Solution to LAP

| | , | 1 | | 2 | | 3 | 4 | 4 | Ţ | 5 |
|---|-------|---------------|----|------------------|----|---------------|----|-----------------------|----|--------------|
| 1 | 0. | 95 | 0. | 7 <i>(</i> 70 | 0. | 62 | 0. | 11 | 0. | 6 |
| 2 | 0. | 23 | 0. | 1 0 | 0. | 70 79 | 0. | 94 | 0. | 35 |
| 3 | 0. | 7 4 J 1 | 0. | 92 | 0. | 92 | 0. |) <u>2</u> | 0. | 31 |
| 4 | 0. | 19 | 0. | 82 | 0. | 74 | 0. | 4.4 1 1 | 0. | 31 |
| 5 | Û. | 39 | 0. | 14 | 0. | 18 | 0. | 39 | 0. | 14 |

- Greedy algorithm
 - Find the largest score
 - Remove scores in same row and column from consideration
 - Repeat
- Result: score = 0.95 + 0.94 + 0.92 + 0.82 + 0.14 = 3.77

Is this the best we can do?



Greedy Solution to LAP

| | 1 | 2 | 3 | 4 | 5 |
|---|------|------|------|------|------|
| 1 | 0.95 | 0.76 | 0.62 | 0.41 | 0.06 |
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Greedy solution score = 3.77

Optimal solution score = 4.26

Discussion

- Greedy method is easy to program, quick to run, and yields "pretty good" solutions in practice.
- But it often does not yield the optimal solution.



Optimal Solution

Hungarian Algorithm

- There is an algorithm called Kuhn-Munkres or "Hungarian" algorithm specifically developed to efficiently solve the linear assignment problem.
- Reduces assignment problem to bipartite graph matching.
- > When starting from an $N \times N$ matrix, it runs in $\mathcal{O}(N^3)$.
- \Rightarrow If you need LAP, you should use it.

In the following

- Look at other algorithms that generalize to multi-frame (>2 frames) problems.
- ⇒ Min-Cost Network Flow



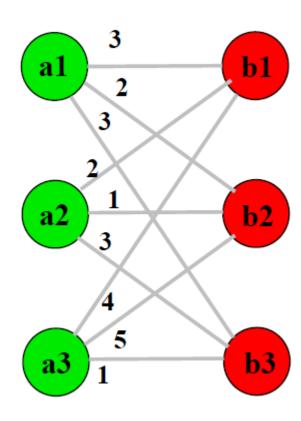
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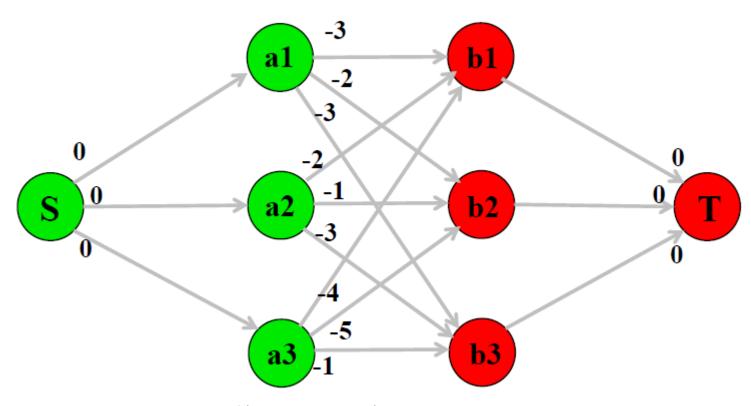


• Small example

| | 1 | 2 | 3 |
|---|---|---|---|
| 1 | 3 | 2 | 3 |
| 2 | 2 | 1 | 3 |
| 3 | 4 | 5 | 1 |

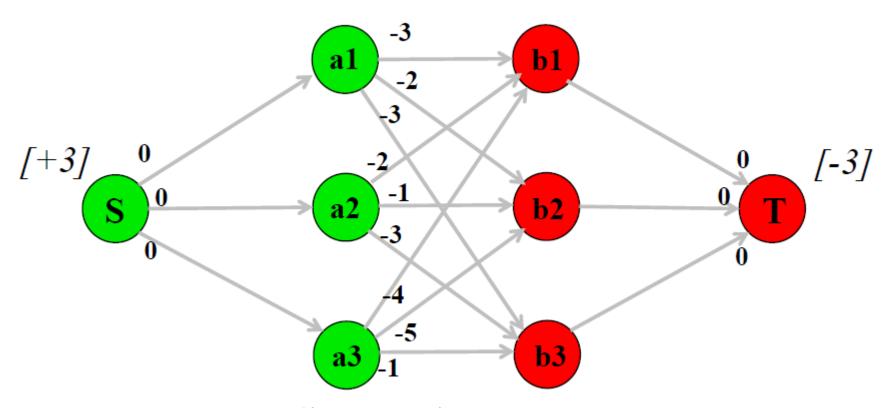






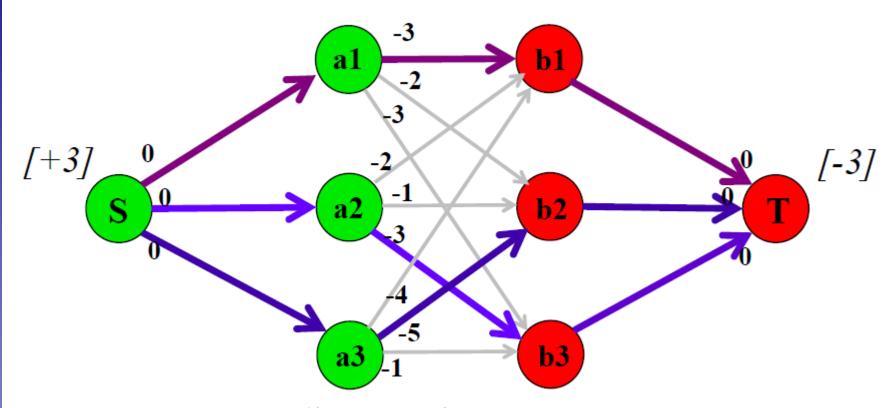
- Conversion into flow graph
 - ightarrow Transform weights into costs $\,c_{ij}=lpha-w_{ij}\,$
 - Add source/sink nodes with 0 cost.
 - \triangleright Directed edges with a capacity of 1.





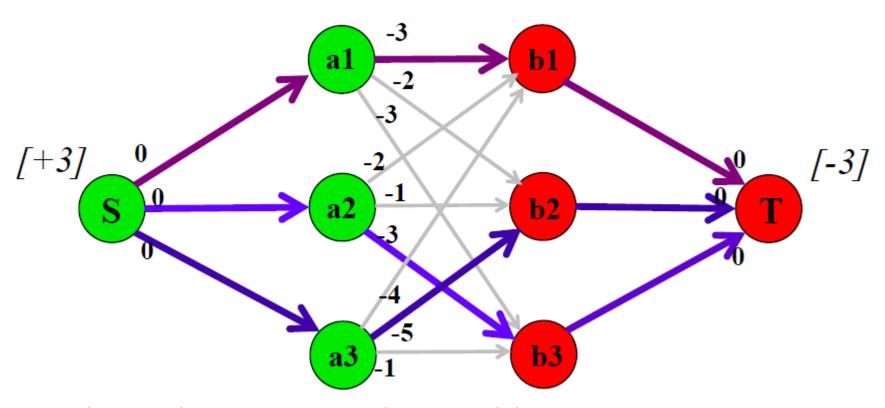
- Conversion into flow graph
 - ightarrow Pump N units of flow from source to sink.
 - Internal nodes pass on flow (Σ flow in $=\Sigma$ flow out).
 - ⇒ Find the optimal paths along which to ship the flow.





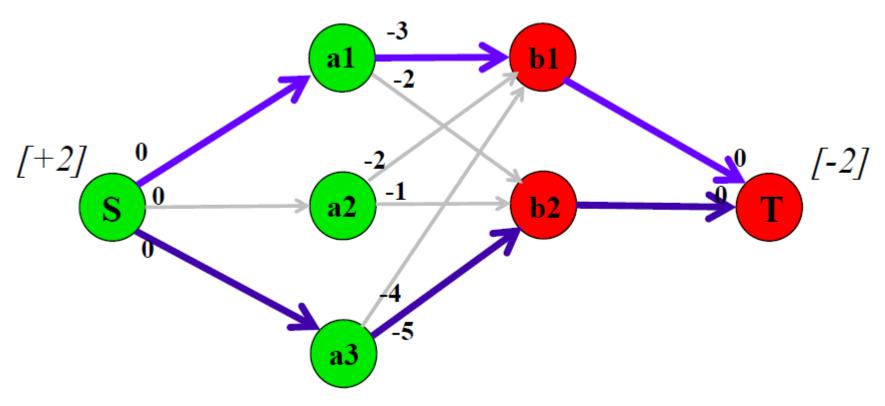
- Conversion into flow graph
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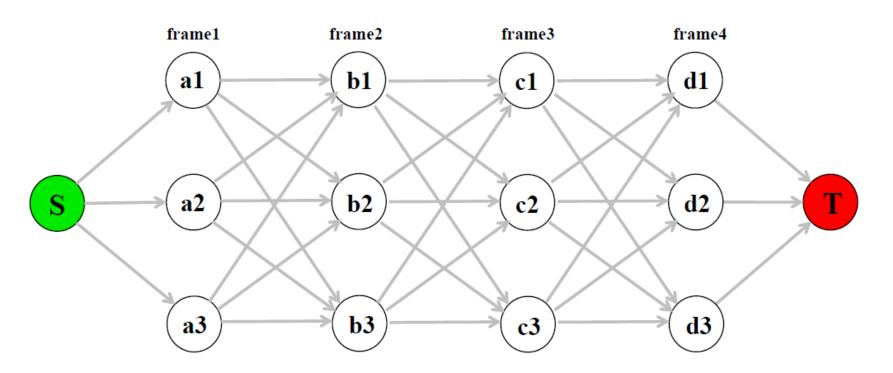
- Solving the Min-Cost Flow problem
 - There are standard algorithms for efficiently solving min-cost network flow
 - E.g., push-relabel or successive shortest path algorithms





- Nice property
 - Min-cost formalism readily generalizes to matching sets with unequal sizes.

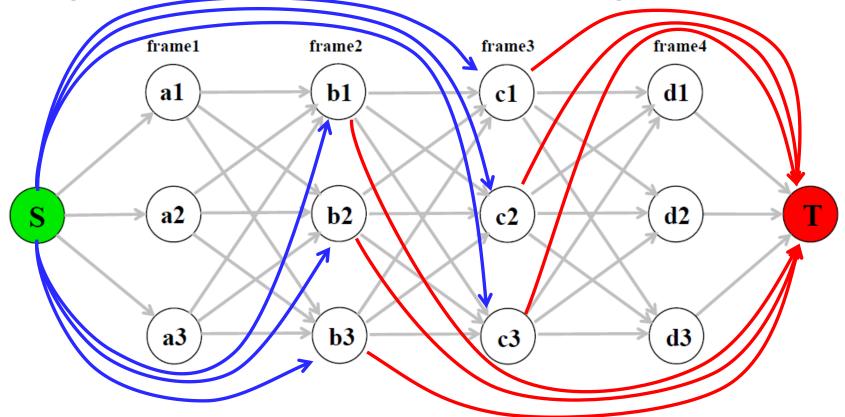




Approach

- Seek a globally optimal solution by considering observations over all frames in "batch mode".
- ⇒ Extend two-frame min-cost formulation by adding observations from all frames into the network.

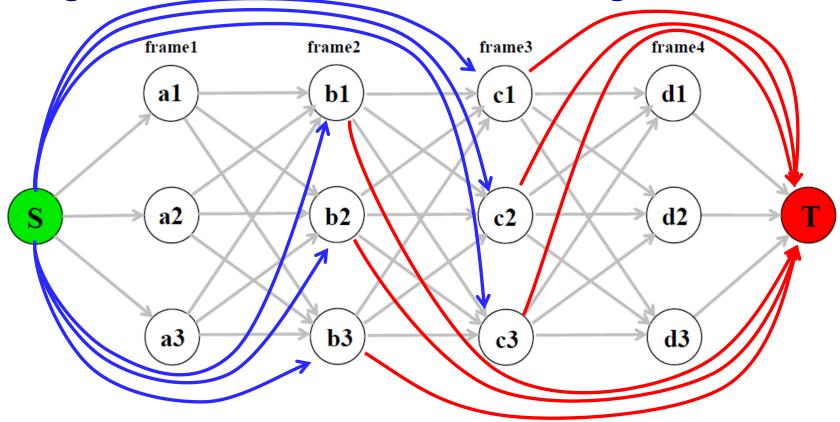




Complication 1

- > Tracks can start later than frame1 (and end earlier than frame4)
- ⇒ Connect the source and sink nodes to all intermediate nodes.

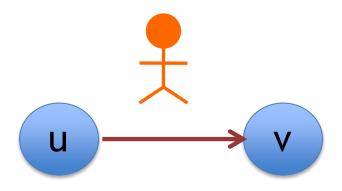




- Complication 2
 - Trivial solution: zero cost!



- Solution
 - Divide each detection into 2 nodes



Detection edge

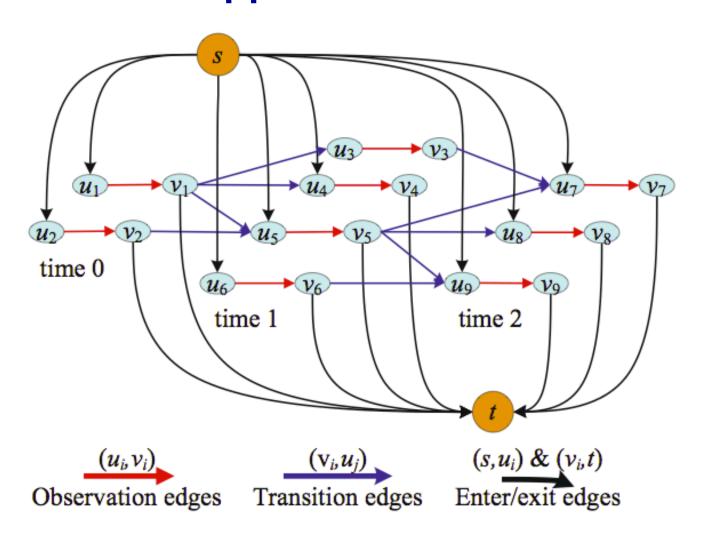
$$C_i = \log \frac{\beta_i}{1 - \beta_i} \longleftarrow$$

Probability that detection i is a false alarm

Zhang, Li, Nevatia, <u>Global Data Association for Multi-Object Tracking using Network Flows</u>, CVPR'08.



Network Flow Approach



Zhang, Li, Nevatia, <u>Global Data Association for Multi-Object Tracking using Network Flows</u>, CVPR'08.



Network Flow Approach: Illustration

Frame t-1

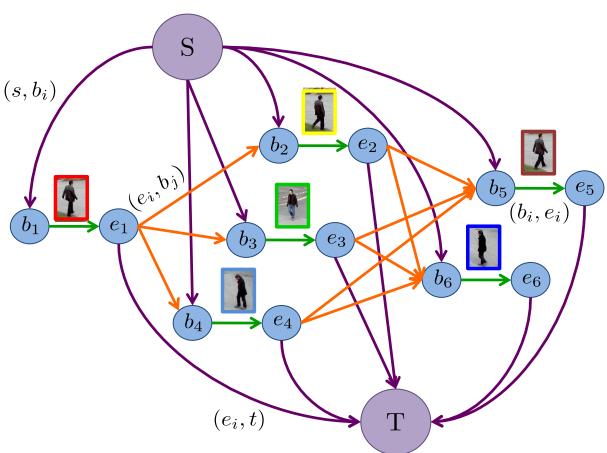


Frame t



Frame t+1







Min-Cost Formulation

Objective Function

$$\mathcal{T}* = \operatorname*{argmin}_{\mathcal{T}} \sum_{i} C_{in,i} f_{in,i} + \sum_{i} C_{i,out} f_{i,out}$$

$$+\sum_{i,j} C_{i,j} f_{i,j} + \sum_{i} C_{i} f_{i}$$

- subject to
 - Flow conservation at all nodes

$$f_{in,i} + \sum_{j} f_{j,i} = f_i = f_{out,i} + \sum_{j} f_{i,j} \ \forall i$$

Edge capacities

$$f_i \leq 1$$



Min-Cost Formulation

Objective Function

$$\mathcal{T}* = \underset{\mathcal{T}}{\operatorname{argmin}} \sum_{i} C_{in,i} f_{in,i} + \sum_{i} C_{i,out} f_{i,out}$$
 $+ \sum_{i,j} C_{i,j} f_{i,j} + \sum_{i} C_{i} f_{i}$ $C_{i} = -log(P_{i})$

Equivalent to Maximum A-Posteriori formulation

$$\mathcal{T}* = \underset{\mathcal{T}}{\operatorname{argmax}} \prod_{i} P(\mathbf{o}_{i}|\mathcal{T})P(\mathcal{T})$$



Min-Cost Formulation

• Objective Function IN OUT $\mathcal{T}* = \mathop{\rm argmin}_{\mathcal{T}} \sum_{i} C_{in,i} f_{in,i} + \sum_{i} C_{i,out} f_{i,out} \\ + \sum_{i} C_{i,j} f_{i,j} + \sum_{i} C_{i} f_{i} \quad \text{Likelihood of the detection} \\ C_{i} = -log(P_{i})$

Equivalent to Maximum A-Posteriori formulation

$$\mathcal{T}* = \underset{\mathcal{T}}{\mathbf{argmax}} \prod_{i} P(\mathbf{o}_i | \mathcal{T}) P(\mathcal{T}) \quad \begin{array}{c} \text{Independence} \\ \text{assumption} \\ P(\mathcal{T}) = \prod_{r_k \in \mathcal{T}} P(T_k) \end{array} \qquad \begin{array}{c} + \\ \text{Markov} \end{array}$$



Network Flow Solutions

Push-relabel method

Zhang, Li and Nevatia, "Global Data Association for Multi-Object Tracking Using Network Flows," CVPR 2008.

Successive shortest path algorithm

- Berclaz, Fleuret, Turetken and Fua, "Multiple Object Tracking using K-shortest Paths Optimization," IEEE PAMI, Sep 2011.
- Pirsiavash, Ramanan, Fowlkes, "Globally Optimal Greedy Algorithms for Tracking a Variable Number of Objects", CVPR'11.
- > These both include approximate dynamic programming solutions



References and Further Reading

- The original network flow tracking paper
 - > Zhang, Li, Nevatia, <u>Global Data Association for Multi-Object</u> <u>Tracking using Network Flows</u>, CVPR'08.
- Extensions and improvements
 - Berclaz, Fleuret, Turetken, Fua, <u>Multiple Object Tracking using</u>
 <u>K-shortest Paths Optimization</u>, IEEE PAMI, Sep 2011. (<u>code</u>)
 - Pirsiavash, Ramanan, Fowlkes, <u>Globally Optimal Greedy</u>
 <u>Algorithms for Tracking a Variable Number of Objects</u>, CVPR'11.
- A recent extension to incorporate social walking models
 - L. Leal-Taixe, G. Pons-Moll, B. Rosenhahn, <u>Everybody Needs</u> <u>Somebody: Modeling Social and Grouping Behavior on a Linear</u> <u>Programming Multiple People Tracker</u>, ICCV Workshops 2011.