

RWTH AACHEN
UNIVERSITY

Computer Vision II - Lecture 13

Multi-Object Tracking III

03.07.2014

Bastian Leibe
RWTH Aachen
<http://www.vision.rwth-aachen.de>
leibe@vision.rwth-aachen.de

Computer Vision II, Summer'14

RWTH AACHEN
UNIVERSITY

Course Outline

- Single-Object Tracking
- Bayesian Filtering
 - Kalman filters
 - Particle filters
 - Case studies
- Multi-Object Tracking
 - Introduction
 - MHT, JPDAF
 - Network Flow Optimization
- Articulated Tracking

Computer Vision II, Summer'14

RWTH AACHEN
UNIVERSITY

Topics of This Lecture

- Recap: MHT
- Data Association as Linear Assignment Problem
 - LAP formulation
 - Greedy algorithm
 - Hungarian algorithm
- Tracking as Network Flow Optimization
 - Min-cost network flow
 - Generalizing to multiple frames
 - Complications
 - Formulation

B. Leibe 3

Computer Vision II, Summer'14

RWTH AACHEN
UNIVERSITY

Recap: Multi-Hypothesis Tracking (MHT)

- Ideas
 - Instead of forming a track tree, keep a set of hypotheses that generate child hypotheses based on the associations.
 - Enforce exclusion constraints between tracks and measurements in the assignment process.
 - Integrate track generation into the assignment process.
 - After hypothesis generation, merge and prune the current hypothesis set.

D. Reid, [An Algorithm for Tracking Multiple Targets](#), IEEE Trans. Automatic Control, Vol. 24(6), pp. 843-854, 1979.
B. Leibe 5

Computer Vision II, Summer'14

Image source: ICox, IJCV'93

RWTH AACHEN
UNIVERSITY

Recap: Hypothesis Generation

- Create hypothesis matrix of the feasible associations

$$\Theta = \begin{bmatrix} 1 & 0 & 1 & 1 \\ 1 & 1 & 1 & 1 \\ 0 & 1 & 1 & 1 \\ 0 & 0 & 1 & 1 \end{bmatrix} \quad \begin{matrix} \mathbf{x}_1 & \mathbf{x}_2 & \mathbf{x}_{fa} & \mathbf{x}_{nt} \\ \mathbf{y}_1 & \mathbf{y}_2 & \mathbf{y}_3 & \mathbf{y}_4 \end{matrix}$$

- Interpretation
 - Columns represent tracked objects, rows encode measurements
 - A non-zero element at matrix position (i,j) denotes that measurement y_i is contained in the validation region of track x_j .
 - Extra column \mathbf{x}_{fa} for association as false alarm.
 - Extra column \mathbf{x}_{nt} for association as new track.
 - Turn this hypothesis matrix

B. Leibe 6

Computer Vision I, Summer'14

RWTH AACHEN
UNIVERSITY

Recap: Assignments

Z_j	\mathbf{x}_1	\mathbf{x}_2	\mathbf{x}_{fa}	\mathbf{x}_{nt}
\mathbf{y}_1	0	0	1	0
\mathbf{y}_2	1	0	0	0
\mathbf{y}_3	0	1	0	0
\mathbf{y}_4	0	0	0	1

- Impose constraints
 - A measurement can originate from only one object.
⇒ Any row has only a single non-zero value.
 - An object can have at most one associated measurement per time step.
⇒ Any column has only a single non-zero value, except for \mathbf{x}_{fa} , \mathbf{x}_{nt}

B. Leibe 7

Computer Vision I, Summer'14

RWTH AACHEN UNIVERSITY

Recap: Calculating Hypothesis Probabilities

- Probabilistic formulation
 - It is straightforward to enumerate all possible assignments.
 - However, we also need to calculate the probability of each child hypothesis.
 - This is done recursively:

$$\begin{aligned}
 p(\Omega_j^{(k)} | \mathbf{Y}^{(k)}) &= p(Z_j^{(k)}, \Omega_{p(j)}^{(k-1)} | \mathbf{Y}^{(k)}) \\
 &\stackrel{\text{Bayes}}{=} \eta p(\mathbf{Y}^{(k)} | Z_j^{(k)}, \Omega_{p(j)}^{(k-1)}) p(Z_j^{(k)}, \Omega_{p(j)}^{(k-1)}) \\
 &= \eta p(\mathbf{Y}^{(k)} | Z_j^{(k)}, \Omega_{p(j)}^{(k-1)}) \underbrace{p(Z_j^{(k)} | \Omega_{p(j)}^{(k-1)})}_{\text{Prob. of assignment set}} \underbrace{p(\Omega_{p(j)}^{(k-1)})}_{\text{Prob. of parent}}
 \end{aligned}$$

Normalization factor Measurement likelihood Prob. of assignment set Prob. of parent

B. Leibe

Computer Vision II, Summer'14

8

RWTH AACHEN UNIVERSITY

Recap: Measurement Likelihood

- Use KF prediction
 - Assume that a measurement $y_i^{(k)}$ associated to a track x_j has a Gaussian pdf centered around the measurement prediction $\hat{x}_j^{(k)}$ with innovation covariance $\hat{\Sigma}_j^{(k)}$.
 - Further assume that the pdf of a measurement belonging to a new track or false alarm is uniform in the observation volume W (the sensor's field-of-view) with probability W^{-1} .
 - Thus, the measurement likelihood can be expressed as

$$\begin{aligned}
 p(\mathbf{Y}^{(k)} | Z_j^{(k)}, \Omega_{p(j)}^{(k-1)}) &= \prod_{i=1}^{M_k} \mathcal{N}(y_i^{(k)}; \hat{x}_j, \hat{\Sigma}_j^{(k)})^{\delta_i} W^{-(1-\delta_i)} \\
 &= W^{-(N_{\text{fal}}+N_{\text{new}})} \prod_{i=1}^{M_k} \mathcal{N}(y_i^{(k)}; \hat{x}_j, \hat{\Sigma}_j^{(k)})^{\delta_i}
 \end{aligned}$$

B. Leibe

Computer Vision II, Summer'14

9

RWTH AACHEN UNIVERSITY

Recap: Probability of an Assignment Set

$$p(Z_j^{(k)} | \Omega_{p(j)}^{(k-1)})$$

- Composed of three terms
 1. Probability of the number of tracks N_{det} , N_{fal} , N_{new}
 - Assumption 1: N_{det} follows a binomial distribution
 2. Probability of a specific assignment of measurements
 - Such that $M_k = N_{\text{det}} + N_{\text{fal}} + N_{\text{new}}$ holds.
 - This is determined as 1 over the number of combinations
 3. Probability of a specific assignment of tracks
 - Given that a track can be either detected or not detected.
 - This is determined as 1 over the number of assignments

$$\begin{aligned}
 p(N_{\text{det}}, N_{\text{fal}}, N_{\text{new}} | \Omega_{p(j)}^{(k-1)}) &= \binom{N}{N_{\text{det}}} p_{\text{det}}^{N_{\text{det}}} (1-p_{\text{det}})^{(N-N_{\text{det}})} \\
 &\quad \text{where } N \text{ is the number of tracks in the parent hypothesis} \\
 &\quad \text{Assumption 2: } N_{\text{fal}} \text{ and } N_{\text{new}} \text{ both follow a Poisson distribution with expected number of events } \lambda_{\text{fal}} W \text{ and } \lambda_{\text{new}} W \\
 p(N_{\text{det}}, N_{\text{fal}}, N_{\text{new}} | \Omega_{p(j)}^{(k-1)}) &= \binom{N}{N_{\text{det}}} p_{\text{det}}^{N_{\text{det}}} (1-p_{\text{det}})^{(N-N_{\text{det}})} \\
 &\quad \cdot \mu(N_{\text{fal}}; \lambda_{\text{fal}} W) \cdot \mu(N_{\text{new}}; \lambda_{\text{new}} W)
 \end{aligned}$$

B. Leibe

Computer Vision III, Summer'14

10

RWTH AACHEN UNIVERSITY

Recap: Probability of an Assignment Set

2. Probability of a specific assignment of measurements
 - Such that $M_k = N_{\text{det}} + N_{\text{fal}} + N_{\text{new}}$ holds.
 - This is determined as 1 over the number of combinations
3. Probability of a specific assignment of tracks
 - Given that a track can be either detected or not detected.
 - This is determined as 1 over the number of assignments

$$\frac{N!}{(N - N_{\text{det}})!} \binom{N - N_{\text{det}}}{N_{\text{det}}}$$

⇒ When combining the different parts, many terms cancel out!

B. Leibe

Computer Vision III, Summer'14

11

RWTH AACHEN UNIVERSITY

Laser-based Leg Tracking using MHT

K. Arras, S. Grzonka, M. Luber, W. Burgard, [Efficient People Tracking in Laser Range Data using a Multi-Hypothesis Leg-Tracker with Adaptive Occlusion Probabilities](#), ICRA'08.

video source: Social Robotics Lab., Univ. Freiburg

Computer Vision I, Summer'14

12

RWTH AACHEN UNIVERSITY

Laser-based People Tracking using MHT

Multi Hypothesis Tracking of People

Matthias Luber, Gian Diego Tipaldi and Kai O. Arras

Laser-based People Tracking using MHT
(Inner city of Freiburg, Germany)
Results projected onto video data.

Social Robotics Laboratory

UNI FREIBURG

video source: Social Robotics Lab., Univ. Freiburg

Computer Vision I, Summer'14

13

Topics of This Lecture

- Recap: MHT
- Data Association as Linear Assignment Problem**
 - LAP formulation
 - Greedy algorithm
 - Hungarian algorithm
- Tracking as Network Flow Optimization
 - Min-cost network flow
 - Generalizing to multiple frames
 - Complications
 - Formulation

Computer Vision II, Summer'14 RWTH AACHEN UNIVERSITY

B. Leibe 14

Back to Data Association...

- Goal:** Match detections across frames

Computer Vision II, Summer'14 RWTH AACHEN UNIVERSITY

Slide credit: Robert Collins B. Leibe 15

Data Association

- Main question here**
 - How to determine which measurements to add to which track?
 - Today: consider this as a matching problem

Computer Vision II, Summer'14 RWTH AACHEN UNIVERSITY

Slide credit: Robert Collins B. Leibe 16

Linear Assignment Formulation

- Form a matrix of pairwise similarity scores
- Similarity could be
 - based on motion prediction
 - based on appearance
 - based on both

		Frame $t+1$	
Frame t		0.11	0.95
		0.85	0.25
		0.90	0.12
		0.81	

Computer Vision II, Summer'14 RWTH AACHEN UNIVERSITY

Slide credit: Robert Collins B. Leibe 17

- Goal**
 - Choose one match from each row and column to maximize the sum of scores

Linear Assignment Formulation

- Example: Similarity based on motion prediction**
 - Predict motion for each trajectory and assign scores for each measurement based on inverse (Mahalanobis) distance, such that closer measurements get higher scores.

Computer Vision I, Summer'14 RWTH AACHEN UNIVERSITY

Slide credit: Robert Collins B. Leibe 18

- Choose at most one match in each row and column to maximize sum of scores

Linear Assignment Problem

- Formal definition**
 - Maximize $\sum_{i=1}^N \sum_{j=1}^M w_{ij} z_{ij}$
 - subject to $\sum_{j=1}^M z_{ij} = 1; i = 1, 2, \dots, N$
 - $\sum_{i=1}^N z_{ij} = 1; j = 1, 2, \dots, M$
 - $z_{ij} \in \{0, 1\}$

Computer Vision I, Summer'14 RWTH AACHEN UNIVERSITY

The constraints ensure that Z is a permutation matrix

- The permutation matrix constraint ensures that we can only match up one object from each row and column.
- Note: Alternatively, we can minimize cost rather than maximizing weights.

$$\arg \min_{z_{ij}} \sum_{i=1}^N \sum_{j=1}^M c_{ij} z_{ij}$$

Slide adapted from Robert Collins B. Leibe 19

Greedy Solution to LAP

	1	2	3	4	5
1	0.95	0.76	0.62	0.41	0.06
2	0.23	0.46	0.79	0.94	0.35
3	0.61	0.02	0.92	0.92	0.81
4	0.49	0.82	0.74	0.41	0.01
5	0.89	0.44	0.18	0.89	0.14

- Greedy algorithm
 - Find the largest score
 - Remove scores in same row and column from consideration
 - Repeat
- Result: score =

Slide credit: Robert Collins B. Leibe 20

Greedy Solution to LAP

	1	2	3	4	5
1	0.95	0.76	0.62	0.41	0.06
2	0.23	0.46	0.79	0.94	0.35
3	0.61	0.02	0.92	0.92	0.81
4	0.49	0.82	0.74	0.41	0.01
5	0.89	0.44	0.18	0.89	0.14

- Greedy algorithm
 - Find the largest score
 - Remove scores in same row and column from consideration
 - Repeat
- Result: score = $0.95 + 0.94 + 0.92 + 0.82 + 0.14 = 3.77$

Is this the best we can do?

Slide credit: Robert Collins B. Leibe 21

Greedy Solution to LAP

	1	2	3	4	5
1	0.95	0.76	0.62	0.41	0.06
2	0.23	0.46	0.79	0.94	0.35
3	0.61	0.02	0.92	0.92	0.81
4	0.49	0.82	0.74	0.41	0.01
5	0.89	0.44	0.18	0.89	0.14

Greedy solution score = 3.77

	1	2	3	4	5
1	0.95	0.76	0.62	0.41	0.06
2	0.23	0.46	0.79	0.94	0.35
3	0.61	0.02	0.92	0.92	0.81
4	0.49	0.82	0.74	0.41	0.01
5	0.89	0.44	0.18	0.89	0.14

Optimal solution score = 4.26

- Discussion
 - Greedy method is easy to program, quick to run, and yields “pretty good” solutions in practice.
 - But it often does not yield the optimal solution.

Slide credit: Robert Collins B. Leibe 22

Optimal Solution

- Hungarian Algorithm
 - There is an algorithm called Kuhn-Munkres or “Hungarian” algorithm specifically developed to efficiently solve the linear assignment problem.
 - Reduces assignment problem to bipartite graph matching.
 - When starting from an $N \times N$ matrix, it runs in $\mathcal{O}(N^3)$.
 - ⇒ If you need LAP, you should use it.
- In the following
 - Look at other algorithms that generalize to multi-frame (> 2 frames) problems.
 - ⇒ Min-Cost Network Flow

Slide credit: Robert Collins B. Leibe 23

Topics of This Lecture

- Recap: MHT
- Data Association as Linear Assignment Problem
 - LAP formulation
 - Greedy algorithm
 - Hungarian algorithm
- Tracking as Network Flow Optimization
 - Min-cost network flow
 - Generalizing to multiple frames
 - Complications
 - Formulation

B. Leibe 24

Min-Cost Flow

- Small example

	1	2	3
1	3	2	3
2	2	1	3
3	4	5	1

Slide credit: Robert Collins B. Leibe 25

Min-Cost Flow

- Conversion into flow graph
 - Transform weights into costs $c_{ij} = \alpha - w_{ij}$
 - Add source/sink nodes with 0 cost.
 - Directed edges with a capacity of 1.

Slide credit: Robert Collins B. Leibe 26

Min-Cost Flow

- Conversion into flow graph
 - Pump N units of flow from source to sink.
 - Internal nodes pass on flow (\sum flow in = \sum flow out).

⇒ Find the optimal paths along which to ship the flow.

Slide credit: Robert Collins B. Leibe 27

Min-Cost Flow

- Conversion into flow graph
 - Pump N units of flow from source to sink.
 - Internal nodes pass on flow (\sum flow in = \sum flow out).

⇒ Find the optimal paths along which to ship the flow.

Slide credit: Robert Collins B. Leibe 28

Min-Cost Flow

- Solving the Min-Cost Flow problem
 - There are standard algorithms for efficiently solving min-cost network flow
 - E.g., push-relabel or successive shortest path algorithms

Slide credit: Robert Collins B. Leibe 29

Min-Cost Flow

- Nice property
 - Min-cost formalism readily generalizes to matching sets with unequal sizes.

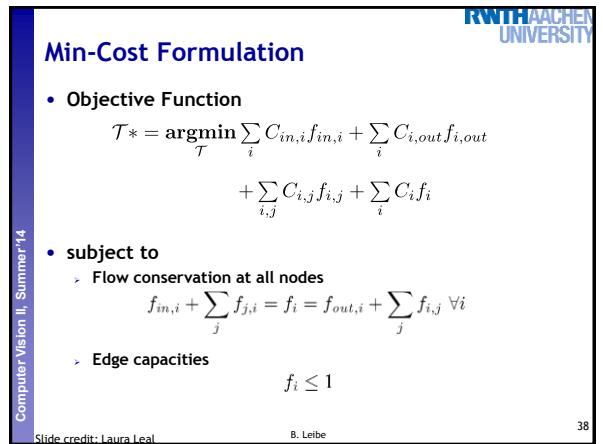
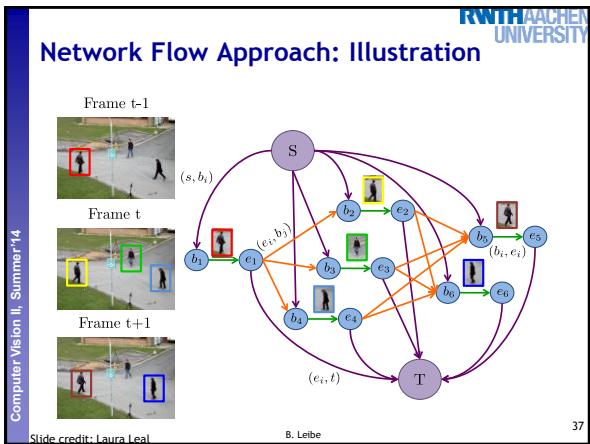
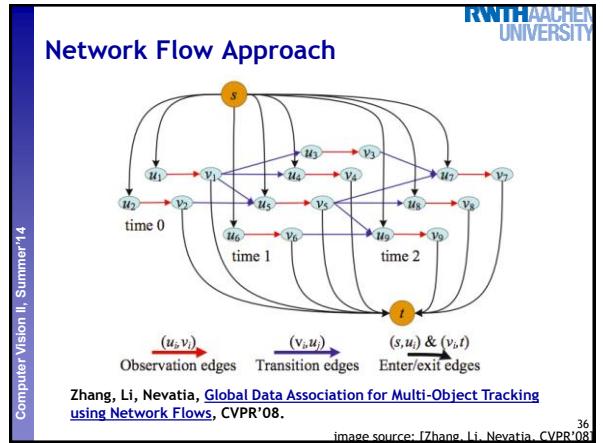
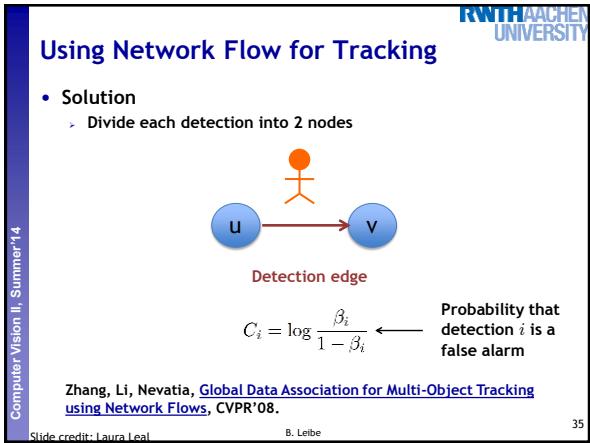
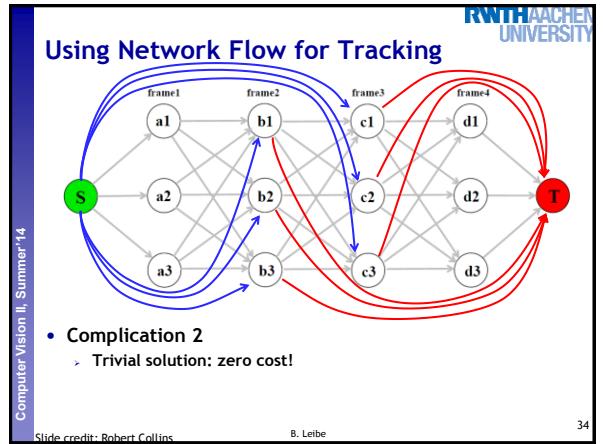
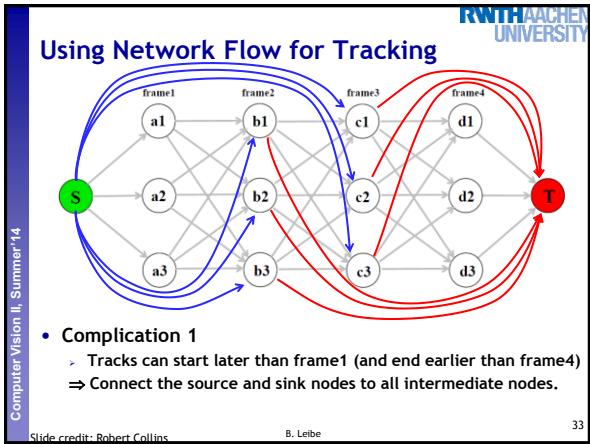
Slide credit: Robert Collins B. Leibe 31

Using Network Flow for Tracking

- Approach
 - Seek a globally optimal solution by considering observations over all frames in “batch mode”.

⇒ Extend two-frame min-cost formulation by adding observations from all frames into the network.

Slide credit: Robert Collins B. Leibe 32



Min-Cost Formulation

• Objective Function

$$\mathcal{T}^* = \operatorname{argmin}_{\mathcal{T}} \sum_i C_{in,i} f_{in,i} + \sum_i C_{i,out} f_{i,out}$$

$$+ \sum_{i,j} C_{i,j} f_{i,j} + \sum_i C_i f_i$$

\downarrow

$C_i = -\log(P_i)$

• Equivalent to Maximum A-Posteriori formulation

$$\mathcal{T}^* = \operatorname{argmax}_{\mathcal{T}} \prod_i P(\mathbf{o}_i|\mathcal{T})P(\mathcal{T})$$

Slide credit: Laura Leal
B. Leibe
39

Min-Cost Formulation

• Objective Function

	IN	OUT	
	$\sum_i C_{in,i} f_{in,i}$	$\sum_i C_{i,out} f_{i,out}$	
	$+ \sum_{i,j} C_{i,j} f_{i,j}$	$+ \sum_i C_i f_i$	Likelihood of the detection

\downarrow

$C_i = -\log(P_i)$

• Equivalent to Maximum A-Posteriori formulation

$$\mathcal{T}^* = \operatorname{argmax}_{\mathcal{T}} \prod_i P(\mathbf{o}_i|\mathcal{T})P(\mathcal{T})$$

$P(\mathcal{T}) = \prod_{T_k \in \mathcal{T}} P(T_k)$

Independence assumption
+ Markov

Slide credit: Laura Leal
B. Leibe
40

Network Flow Solutions

- Push-relabel method
 - Zhang, Li and Nevatia, “Global Data Association for Multi-Object Tracking Using Network Flows,” CVPR 2008.
- Successive shortest path algorithm
 - Berclaz, Fleuret, Turetken and Fua, “Multiple Object Tracking using K-shortest Paths Optimization,” IEEE PAMI, Sep 2011.
 - Pirsiavash, Ramanan, Fowlkes, “Globally Optimal Greedy Algorithms for Tracking a Variable Number of Objects”, CVPR’11.
 - These both include approximate dynamic programming solutions

Slide credit: Robert Collins
B. Leibe
41

References and Further Reading

- The original network flow tracking paper
 - Zhang, Li, Nevatia, [Global Data Association for Multi-Object Tracking using Network Flows](#), CVPR’08.
- Extensions and improvements
 - Berclaz, Fleuret, Turetken, Fua, [Multiple Object Tracking using K-shortest Paths Optimization](#), IEEE PAMI, Sep 2011. ([code](#))
 - Pirsiavash, Ramanan, Fowlkes, [Globally Optimal Greedy Algorithms for Tracking a Variable Number of Objects](#), CVPR’11.
- A recent extension to incorporate social walking models
 - L. Leal-Taixe, G. Pons-Moll, B. Rosenhahn, [Everybody Needs Somebody: Modeling Social and Grouping Behavior on a Linear Programming Multiple People Tracker](#), ICCV Workshops 2011.

Slide credit: Robert Collins
B. Leibe
42